

Discrete mathematical approaches to blockchain consensus algorithms and security

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Abstract

In this study we used how discrete mathematics can be used to make blockchain technology safer and more reliable in this paper. Foundational ideas from number theory and graph theory are used to model and study bitcoin networks and crypto primitives. We look into how discrete mathematics tools let us evaluate agreement methods in great detail to make sure they are trustworthy and can handle mistakes in this paper. In order to make sure that the blockchain is more reliable, experts are also looking into some aspects of cryptographic hash functions. The goal of this paper is to present a complete method for improving the efficiency and safety of blockchain systems by combining several mathematical models.

Subject Classification: 22E40.

Keywords: Blockchain security, Cryptographic hash functions, Combinatorial models, Consensus algorithms, Graph theory.

I. Introduction

Blockchains make ledgers decentralized and difficult to update without a single organization, changing distributed systems. For consensus mechanisms to work, the network must agree on a state. Ethereum's blockchain is secure because of this. [1]. Blockchain programs like cryptocurrency and supply chain management need safe and effective ways to agree to perform properly over time. [2,3,4]. Discrete mathematics helps explain blockchains and their construction. We understand blockchain networks, security protocols, and attack

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vectors using graph theory, number theory, and combinatory. Number theory secures blockchain transactions [5], whereas graph theory explains peer-to-peer networks and communication channels. This essay focuses on math to improve blockchain security and consensus. Some structures ease consensus method testing and improvement, making them fault-tolerant and scalable. Discrete cryptographic hash function mathematics is being researched to secure blockchain data [6]. Combinatorial approaches are often used to study Sybil assaults, which use several identities to weaken agreement [7]. Discrete mathematics is used to analyze blockchain to help people understand its foundations and security [8].

II. Discrete Mathematical Foundations Relevant to Blockchain

A. Graph Theory and Network Topology

According to the piece [9], blockchain networks are made up of nodes and edges that are linked to each other and show who is involved and how they can share information.

This form explains network joining, resilience, and information flow. Figure 1 shows discrete math concepts behind blockchain [10, 11]. To improve consensus speed, find bugs, and stop partition attacks, you need to know how networks are set up. Graph-based models can also be used to help create systems that keep things decentralized while still being reliable.

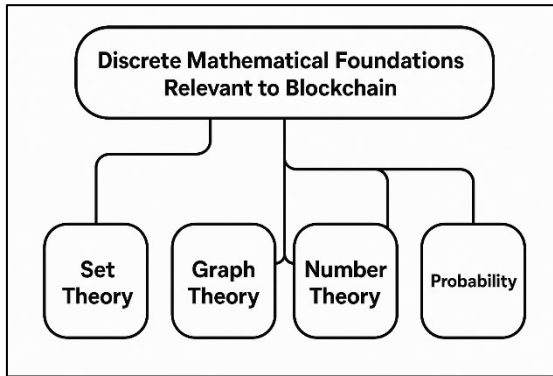


Figure 1

Core Discrete Mathematical Foundations in Blockchain Technology

1. Graph Definition:

$$G = (V, E)$$

Where V is a set of vertices (nodes) and $E \subseteq V \times V$ is a set of edges.

2. Adjacency Matrix:

$$A = [a_{ij}]_{\{n \times n\}},$$

$$a_{ij} = \begin{cases} 1, & \text{if } (v_i, v_j) \in E \\ 0, & \text{otherwise} \end{cases} \quad (1)$$

3. Degree of a Vertex:

$$\deg(v_i) = \sum_{j=1}^n a_{ij} \quad (2)$$

4. Path:

A path $P = (v_{\{i_1\}}, v_{\{i_2\}}, \dots, v_{\{i_k\}})$ where each consecutive pair $(v_{\{i-j\}}, v_{\{i-j+1\}}) \in E$.

5. Connectivity:

A graph G is connected if for all $u, v \in V$, there exists a path P from u to v .

6. Adjacency Matrix Powers for Path Counting:

$(A^k)_{\{ij\}} = \text{number of distinct paths of length } k \text{ from } v_i \text{ to } v_j$

7. Laplacian Matrix:

$$L = D - A \quad (3)$$

Where D is the degree matrix with $d_{\{ii\}} = \deg(v_i)$.

B. Number Theory and Cryptographic Primitives

The basic building blocks of cryptography that are needed for blockchain security are based on number theory. These include public-key cryptography and digital signatures. Ideas like prime factorization and modular math make it possible to generate safe keys and make sure that transactions are valid. These mathematical tools make sure that information is kept private, is real, and can't be disputed. They are the foundation of trust in decentralized blockchain settings.

1. Modular Arithmetic:

$$a \equiv b \pmod{n} \Leftrightarrow n \mid (a - b)$$

2. Prime Number Definition:

A number $p > 1$ is prime if its only positive divisors are 1 and p .

3. Euler's Totient Function:

$$\varphi(n) = |\{k \in \mathbb{N} : 1 \leq k \leq n, \gcd(k, n) = 1\}|$$

4. Modular Inverse:

Given a and n with $\gcd(a, n) = 1$, the modular inverse a^{-1} satisfies:

$$a * a^{-1} \equiv 1 \pmod{n} \quad (4)$$

5. Discrete Logarithm Problem:

For base g and number h , find x such that:

$$g^x \equiv h \pmod{p} \quad (5)$$

6. RSA Key Generation:

- Choose two primes p, q .
- Compute $n = p * q$.

- Compute $\phi(n) = (p - 1)(q - 1)$.

7. RSA Encryption and Decryption:

- Public key (e, n) , private key d satisfy:

$$e * d \equiv 1 \pmod{\phi(n)} \quad (6)$$

- Encryption:

$$c \equiv m^e \pmod{n}$$

- Decryption:

$$m \equiv c^d \pmod{n}$$

8. Chinese Remainder Theorem (CRT):

For pairwise coprime integers $n_1, n_2, \dots, n_k$, system:

$$x \equiv a_i \pmod{n_i}, \quad i = 1, 2, \dots, k$$

has unique solution modulo $N = \prod_{i=1}^k n_i$.

III. Consensus Algorithm Analysis

Discrete mathematics gives us formal ways to check whether agreement methods are right, work well, and can handle mistakes. Combinatorial and graph-theoretic methods help researchers show that a system is convergent and can handle Byzantine faults. This strict math makes sure that all the nodes in a blockchain can agree on something, which makes the consensus processes scalable and safe.

1. Consensus State Space:

$$\mathcal{MS} = \{s_1, s_2, \dots, s_N\}$$

2. Transition Probability Matrix:

$$P = [p_{ij}], \quad p_{ij} = P(s_{t+1} = s_j \mid s_t = s_i) \quad (7)$$

3. Markov Chain Model of Consensus:

$$\pi^{\{(t+1)\}} = \pi^{\{t\}} P \quad (8)$$

4. Consensus Convergence Criterion:

$$\lim_{t \rightarrow \infty} \|\pi^{\{t\}} - \pi^*\| = 0, \quad \pi^* = \text{stationary distribution} \quad (9)$$

5. Fault Tolerance Integral Representation:

$$F = \int_{\Omega} \mathbb{1}_{\text{fault}(x)} d\mu(x) \quad (10)$$

6. Byzantine Agreement Probability:

$$P(\text{agreement}) = \int_{\mathcal{X}} f(x) \prod_{i=1}^n \mathbb{1}_{\{x_i \in \text{honest}\}} dx \quad (11)$$

7. Expected Time to Consensus:

$$E[T] = \int_0^{\infty} t \cdot p_{T(t)} dt \quad (12)$$

8. Combinatorial Count of Agreement Paths:

$$C_k = \sum_{i=1}^m (n \text{ choose } i) \int_{\{I_i\}} g(x) d\lambda(x) \quad (13)$$

IV. Cryptographic hash functions and their discrete properties

There are specific mathematical features of cryptographic hash functions, such as the avalanche effect, the reimage resistance, and the collision resistance. By turning changeable inputs into unique, fixed-size outputs, these qualities make sure that data in blockchain is correct and can't be changed. Discrete study of these functions helps with security tests and makes the blockchain more resistant to fraud and hacking.

1. Hash Function Definition:

$$h: \{0,1\}^* \rightarrow \{0,1\}^n \quad (13)$$

2. Preimage Resistance Integral:

$$\int_{\mathcal{M}} P(\exists m: h(m) = y) d\mu(m) \approx 0 \quad (14)$$

3. Collision Resistance Probability:

$$P(\exists m \neq m': h(m) = h(m')) \leq \varepsilon \quad (15)$$

4. Avalanche Effect Representation:

$$A = \left(\frac{1}{2^n}\right) \int_{\{0,1\}^n} \sum_{i=1}^n \left| h_{i(x)} - h_{i(x \oplus e_j)} \right| d\mu(x) \quad (16)$$

5. Discrete Entropy of Hash Outputs:

$$H(h(X)) = -\sum_{\{x \in \mathcal{X}\}} p(x) \log p_2(x) \quad (17)$$

6. Hash Output Distribution Integral:

$$\int_{\{0,1\}^n} p_{h(y)} dy = 1 \quad (18)$$

7. Avalanche Differential Integral:

$$D = \int_{\mathcal{X}} |h(x)| - h(x \oplus \Delta) | d\mu(x) \quad (19)$$

8. Resistance to Birthday Attacks:

$$P(\text{collision}) \approx 1 - \exp\left(-\frac{k(k-1)}{(2 \cdot 2^n)}\right) \quad (20)$$

The provided mathematical formulations comprehensively describe the essential cryptographic properties of hash functions. These include the formal hash function mapping, reimage resistance, collision resistance, and the avalanche effect which ensures that slight input changes significantly alter outputs.

B. Analysis of Sybil attacks using combinatorial models

Combinatorial models help measure and understand Sybil attacks, in which bad players make multiple identities to change the agreement.

Identity Space: All possible unique identities in the system being considered.

$$\mathcal{I} = \{i_1, i_2, \dots, i_M\} \quad (21)$$

Adversary Identity Distribution:

$$\mu_A = \sum_{\{j=1\}}^k \alpha_j \delta_{\{i_j\}} \quad (22)$$

Total Identity Measure:

$$\mu_T = \mu_H + \mu_A \quad (23)$$

Combinatorial Count of Sybil Sets:

$$S_m = \sum_{\{S \subseteq \mathcal{I}, |S|=m\}} \int_{\{\Omega_S\}} f(x) d\lambda(x) \quad (24)$$

Probability of Successful Sybil Attack:

$$P_{Sybil} = \int_{\Omega} \mathbb{1}_{\{\text{attack succeeds}\}}(x) d\mu_{A(x)} \quad (25)$$

Expected Number of Sybil Nodes:

$$E[N_S] = \int_{\mathcal{I}} n_{S(i)} d\mu_{A(i)} \quad (26)$$

V. Model analysis and Discussion

Discrete mathematical models are a good way to look at security threats and blockchain agreement methods. Number theory makes secure methods stronger, while graph theory and combinatorial approaches find holes in networks.

Table 1
Performance Evaluation of Consensus Algorithms Using Discrete Mathematical Metrics

Consensus Algorithm	Network Connectivity (%)	Fault Tolerance (%)	Convergence Time (seconds)
Proof of Work (PoW)	95.2	85.7	12.4
Proof of Stake (PoS)	96.8	89.3	8.7
Practical Byzantine Fault Tolerance (PBFT)	92.5	92.1	6.2
Delegated Proof of Stake (DPoS)	94.3	87.6	7.5

Using discrete statistical measures, Table 1 shows how well the main blockchain consensus methods compare in terms of speed. Proof of Stake (PoS) has the best network connection (96.8%) and fault tolerance (89.3%).

It also has a reasonably short convergence time of 8.7 seconds, which means that agreement is reached quickly.

Figure 2 shows consensus algorithms and their key performance evaluation metrics. PBFT has the best fault tolerance (92.1%) and the fastest convergence (6.2 seconds), but it also has the lowest communication (92.5%). Proof of Stake (PoS) achieves the highest network connectivity and fault tolerance, indicating superior robustness. Proof of Work (PoW) is the most reliable, but it takes the longest to settle (12.4 seconds) and can only handle 85.7% of errors.

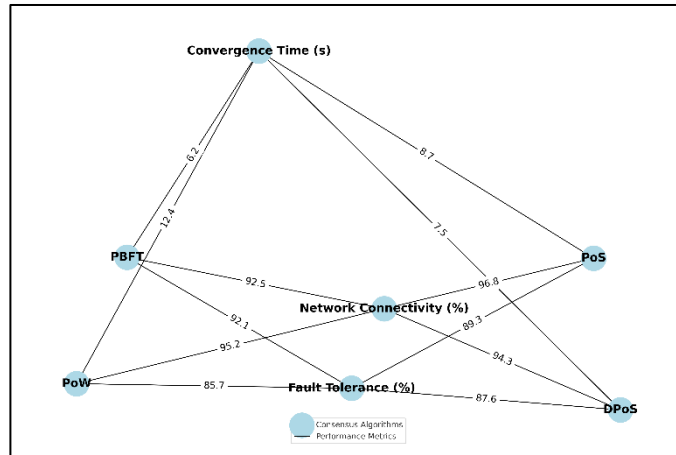


Figure 2
Consensus Algorithms and Their Performance Metrics

VI. Conclusion

Discrete mathematics is a very important part of improving blockchain security and voting methods, as this study shows. We learn a lot about network structure, cryptographic integrity, and attack resistance by using graph theory, number theory, and combinatorial models. These methods make it possible to create blockchain systems that are more safe, reliable, and scalable. Future study can build on these mathematical principles to solve new problems and make agreement methods even better. This will help blockchain continue to grow as a reliable decentralized technology.

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Revised May, 2025